

Lecture 16 - Nov 10

Graphs

Dijkstra's Algorithm: Tracing
Dijkstra's Algorithm: Pre- and Post-cond.

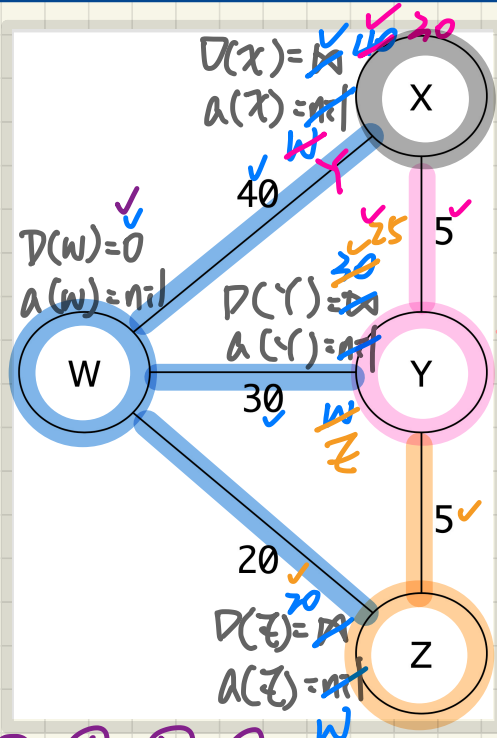
Announcements/Reminders

- Today's class: notes template posted
- **Test 1** results released on Tuesday (Nov 4)
- Change of Dates:
 - + **Assignment 2** to be released on Wed, Nov 12
 - + **Assignment 2** to be due on Wed, Nov 19
 - + **Test 2** to be take place on Mon, Nov 24

Dijkstra's Shortest Path Algorithm: Example 2

single source.

* # Iterations = |V|



$d(w, w) = 0$
 $d(w, x) = 30$
 $d(w, y) = 25$
 $d(w, z) = 20$

Int.

	W	X	Y	Z
∇	0	∞	∞	∞
a	nil	nil	nil	nil

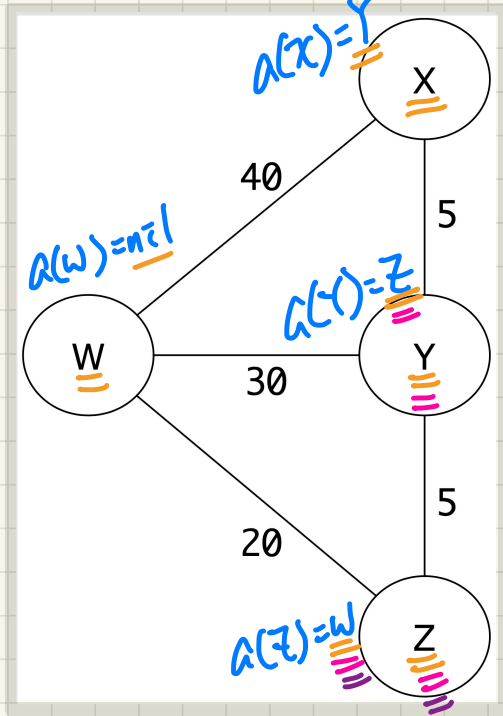
 $\nabla(x) = d(w, x)$

# Iteration	min ∇	changes
1	W $\nabla(w)=0$	x, y, z in Q $a(x)=w \leftarrow 0 + 40 < \infty \rightarrow \nabla(x)=40 \neq d(w, x)$ $a(y)=w \leftarrow 0 + 30 < \infty \rightarrow \nabla(y)=30 \neq d(w, y)$ $a(z)=w \leftarrow 0 + 20 < \infty \rightarrow \nabla(z)=20 = d(w, z)$ (not final) remove W from Q. $\hookrightarrow \nabla(w) = d(w, w)$ final
2	Z $\nabla(z)=20$	$w \notin Q \rightarrow$ not considered. $a(y)=z \leftarrow 20 + 5 < 30 \rightarrow \nabla(y)=25 = d(w, y)$ not final remove Z from Q $\nabla(z) = d(w, z)$ final
3	Y $\nabla(y)=25$	$w \notin Q \rightarrow$ not considered. $a(x)=y \leftarrow 25 + 5 < 40 \rightarrow \nabla(x)=30 = d(w, x)$ not final remove Y from Q $\nabla(y) = d(w, y)$ final
4	X $\nabla(x)=30$	$w \notin Q \wedge y \notin Q \rightarrow$ not considered. **

Θ^*
 ① ④ ③ ②
~~W~~ ~~X~~ ~~Y~~ ~~Z~~

* min key (D value) priority queue

Upon termination of Dijkstra's algorithm

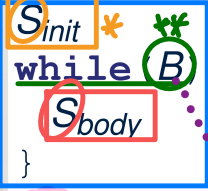


dest.	ancestor path.			
<u>X</u>	<u>X</u>	<u>Y</u>	<u>Z</u>	<u>W</u>
<u>Y</u>		<u>Y</u>	<u>Z</u>	<u>W</u>
<u>Z</u>			<u>Z</u>	<u>W</u>

dest.	shortest path (from source W)			
<u>X</u>	<u>W</u>	<u>Z</u>	<u>Y</u>	<u>X</u>
<u>Y</u>	<u>W</u>	<u>Z</u>	<u>Y</u>	
<u>Z</u>	<u>W</u>	<u>Z</u>		

Correctness of Loops: Syntax

precondition



postcondition

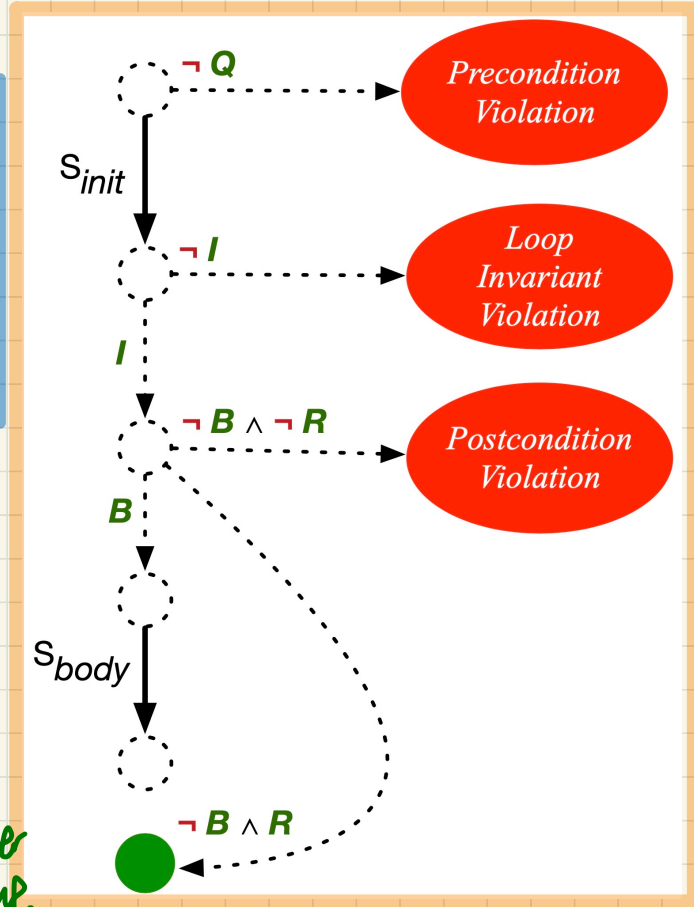
```
void myAlgorithm() {  
  assert Q; /* Precondition */  
  Sinit  
  assert I; /* Is LI established? */  
  while ( B ) {  
    Sbody  
    assert I; /* Is LI preserved? */  
  }  
  assert R; /* Postcondition */  
}
```

Iterative Algorithm

B: stay condition
¬B: exit condition

* Initialization/Preparation for iterations.

** As long as B is true, execute S_{body} another time.
As soon as B is false, exit from the loop.



$\{Q\}$ → precondition:
 $\underline{w(u,v) \geq 0 \quad u \in V \quad v \in V}$
 non-negative weight

Dijkstra's algorithm

Q. when is this $D(u)$ finalized?
 A. when u is removed from Q

shortest-path length known so far

$\{R\}$

implicitly shortest path so far from start vertex x to v
 true shortest-path length between x and v

$D(v) = d(x, v)$

where x is the start vertex
 2. Reverse of ancestor path gives the shortest path.